

The inner model $C(aa^+)$

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Extended constructibility

Research program started by Kennedy, Magidor, Väänänen ([2], [3]).

Definition

Suppose $\mathcal{L}^*$ is a logic. The hierarchy (J'_α) , α a limit ordinal, of sets *construtable using $\mathcal{L}^*$* and the class Tr are defined by transfinite double induction, as follows:

$$\text{Tr} = \{(\alpha, \varphi(\vec{a})) : (J'_\alpha, \in, \text{Tr} \upharpoonright \alpha) \models \varphi(\vec{a}), \varphi(\vec{x}) \in \mathcal{L}^*, \vec{a} \in J'_\alpha, \alpha \in \text{Lim}\},$$

where

$$\text{Tr} \upharpoonright \alpha = \{(\beta, \vec{a}) \in \text{Tr} : \beta < \alpha\},$$

and

$$\begin{aligned} J'_0 &= \emptyset, \\ J'_{\alpha+\omega} &= \text{rud}_{\text{Tr}}(J'_\alpha \cup \{J'_\alpha\}), \\ J'_{\omega\delta} &= \bigcup_{\alpha < \delta} J'_{\omega\alpha} \text{ for limit } \delta. \end{aligned}$$

The class $\bigcup_{\alpha \in \text{Ord}} J'_\alpha$ is denoted by $C(\mathcal{L}^*)$.

Inner model $C(aa)$

- $C(aa)$ obtained from stationary logic
- Stationary logic $\mathcal{L}(aa)$ adds to first-order logic the aa -quantifier, defined as follows:

$$M \models aa s \varphi(s, \vec{a})$$

if there is a closed unbounded set $C \subset \mathcal{P}_{\omega_1}(M)$ such that for any $s \in C$,

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- The aa-quantifier can express, e.g., that a set is countable, that an ordinal has countable cofinality, or that a linear order is $\aleph_1$ -like.

The aa^+ -quantifier

- Variant of the aa-quantifier
- Defined intuitively by

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is there is a club $C \subset \mathcal{P}_{\omega_1}(M)$ such that for each $s \in C$

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- For nested aa^+ -quantifiers, the clubs are always taken in M not in the next admissibles
- Advantage over the aa -quantifier: the aa^+ -quantifier can talk about transitive collapses of well-founded sets in M
- $\mathcal{L}(aa^+)$ logic which adds the aa^+ -quantifier to first order logic
- $C(aa^+)$ the inner model obtained from $\mathcal{L}(aa^+)$

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- For transitive A , the next admissible set $A^+ = L_\alpha(A)$, where α is the least such that $(L_\alpha(A), \in) \models KP$.

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- For transitive A , the next admissible set $A^+ = L_\alpha(A)$, where α is the least such that $(L_\alpha(A), \in) \models KP$.
- For a structure $\mathcal{M} = (M, R_i)_{i \in I}$ such that M is not a transitive set or $\in$ is not in the vocabulary of $\mathcal{M}$, a more natural notion of next admissible is based on structures which have the elements of the domain M as urelements
- For $\mathcal{M}$ as above, a set $(\mathcal{M}; A, \in)$ is admissible above $\mathcal{M}$ if $(\mathcal{M}; A, \in) \models KPU$ (KP with urelements) and $M \in A$
- For $\mathcal{M}$ as above, the next admissible of $\mathcal{M}$ is $(\mathcal{M}; A, \in)$ where $A = \bigcap \{B : (\mathcal{M}; B, \in) \text{ admissible above } \mathcal{M}\}$.

Club determinacy

Idea: no definable stationary co-stationary sets.

Definition

The inner model $C(aa)$ is said to be Club Determined if for all α and for all $\varphi(\vec{x}, \vec{t}, s) \in \mathcal{L}(aa)$, and for all finite sequences $\vec{t}$ of countable subsets of J'_α :

$$(J'_\alpha, \in, \text{Tr} \upharpoonright \alpha) \models \forall \vec{x} [aa s \varphi(\vec{x}, \vec{t}, s) \vee aa s \neg \varphi(\vec{x}, \vec{t}, s)].$$

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Club Determinacy

Theorems (Kennedy, Magidor, Väänänen)

- *If there is a proper class of Woodin cardinals, then Club Determinacy holds in $C(aa)$.*
- *Suppose $C(aa)$ satisfies Club Determinacy. Then every regular $\kappa \geq \omega_1^V$ is measurable in $C(aa)$.*
- *Suppose $C(aa)$ satisfies Club Determinacy. Then the first-order theory of $C(aa)$ is set forcing absolute.*

Lemma

All the above results hold for $C(aa^+)$ with the same proof.

aa-mice by Kennedy-Magidor-Väänänen

- An aa-mouse is a structure of the form $(J_\alpha^T, \in, T, T^*, (P)_\xi)$.
- For limit $\beta < \alpha$, $T_\beta = \{\varphi(\vec{a}) : (\beta, \vec{a}) \in T\}$ is a complete consistent $\mathcal{L}(\text{aa})$ -theory with parameters from J_β^T that extends the first-order theory of $(J_\beta^T, \in, T \upharpoonright \beta)$.
- T^* is a complete consistent $\mathcal{L}(\text{aa})$ -theory that extends the first-order theory of $(J_\alpha^T, \in, T)$.

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- T^* is a complete consistent $\mathcal{L}(\text{aa})$ -theory that extends the first-order theory of $(J_\alpha^T, \in, T)$.
- By starting from a countable aa mouse $(M_o, \in, T, T^*)$ and iterating the aa-ultrapower construction ω_1 -many times, the images of the previous iterates $\{j_{\alpha\omega_1}[M_\alpha] : \alpha < \omega_1\}$ form a club in M_{ω_1} .
- This allows one to prove that the predicate $T_{\omega_1}^*$ of the ω_1 -iterate is correct about $\mathcal{L}(\text{aa})$ -truth.
- Consequently, the ω_1 -iterate is a level of the $C(\text{aa})$ -hierarchy.

Theorem (Kennedy, Magidor, Väänänen)

If Club Determinacy holds in $C(aa)$, then $C(aa)$ satisfies the Continuum Hypothesis.

Theorem (Goldberg, Steel [1])

Is Club Determinacy holds in $C(aa)$, then $C(aa)$ satisfies the Ultrapower Axiom.

Theorem (Goldberg, Steel [1])

Is Club Determinacy holds in $C(aa)$, then $C(aa)$ satisfies GCH.

Good formulas

- In stationary logic, $\varphi \rightarrow aa s\varphi$, where s does not occur in φ , is valid.
- In $\mathcal{L}(aa^+)$, $\varphi \rightarrow aa^+ s\varphi$ is not valid in general.
- Consider, e.g., a model M such that axiom φ of KP fails in M . Then $M \models \neg\varphi$, but necessarily $M \models aa^+ s\varphi$.

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Then $M \models \neg\varphi$, but necessarily $M \models aa^+ s\varphi$.
- We want to be able to use $\varphi \rightarrow aa^+ s\varphi$ in the definition of aa^+ -mice

Good formulas

- Let $x \in J_{\sup T}^T$ be a shorthand for a first-order formula that says in any model containing J_α^T , where $T \subset \text{Lim} \cap \alpha \times J_\alpha^T$, as a subset that $x \in J_\alpha^T$.

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- For ordinals $\eta_1, \dots, \eta_n$, we let $\Psi(y, \gamma, J_{\sup T}^T, \eta_1, \dots, \eta_n, s_1, \dots, s_k)$ be a shorthand for the formula

$$\begin{aligned} & \exists z (\forall w (w \in z \leftrightarrow w \in J_{\sup T}^T)) \\ & \wedge y \in J_\gamma(w, \in, T, \mathbf{P}_{\eta_1}, \dots, \mathbf{P}_{\eta_n}, s_1, \dots, s_k) \\ & \wedge \neg \exists \gamma' \leq \gamma [J_{\gamma'}(w, \in, T, \mathbf{P}_{\eta_1}, \dots, \mathbf{P}_{\eta_n}, s_1, \dots, s_k) \models KP]. \end{aligned}$$

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- I.e., $\Psi(y, \gamma, J_{\sup T}^T, \eta_1, \dots, \eta_n, s_1, \dots, s_k)$ says in any model containing the next admissible $(J_\alpha^T, \in, T, P_{\eta_1}, \dots, P_{\eta_n}, s_1, \dots, s_k)^+$ that y is in the next admissible of $(J_\alpha^T, \in, T, P_{\eta_1}, \dots, P_{\eta_n}, s_1, \dots, s_k)^+$.

Good formulas

Definition (Good formula)

Suppose φ is an $\mathcal{L}(\text{aa}^+)$ -formula in vocabulary τ_ξ^- . We say that a set $\{\theta' : \theta \text{ a subformula of } \varphi\}$ is an existential specification of φ in vocabulary τ_ξ^- if it satisfies:

1. If θ is atomic, then $\theta' = \theta$.
2. If $\theta = \psi \wedge \gamma$, then $\theta' = \psi' \wedge \gamma'$. If $\theta = \neg\psi$, then $\theta' = \neg\psi'$. If $\theta = \text{aa}^+ s \psi$, then $\theta' = \text{aa}^+ s \psi'$.

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3. If θ is $\exists x\psi(x)$, then θ' is $\exists x(\psi^* \wedge \psi'(x))$ where ψ^* is one of the following:
 - 3.1 $x \in J_{\sup T}^T$.
 - 3.2 There are $\{\eta_1, \dots, \eta_n\} \subset \xi$ and $\{k_1, \dots, k_l\} \subset \omega$, at least one of them nonempty, such that
 - 3.2.1 if $\{k_1, \dots, k_l\} \neq \emptyset$, then for some m , θ is in the scope of $\mathbf{aa}^+s_1, \dots, \mathbf{aa}^+s_m$ in φ , $\{k_1, \dots, k_l\} \subset m$,
 - 3.2.2 ψ^* is $\exists \beta \in \text{Ord}(\Psi(x, \beta, J_{\sup T}^T, \eta_1, \dots, \eta_n, s_{k_1}, \dots, s_{k_l}))$.

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We say that φ is *good* (in τ_ξ^-) if there is a formula ψ (in τ_ξ^-) and an existential specification $\{\theta' : \theta \text{ a subformula of } \psi\}$ of ψ such that $\varphi = \psi'$.

$\mathbf{aa}^+$ -premouse

An $\mathbf{aa}^+$ -premouse in vocabulary $\tau_\xi = \{\in, R_T, R_{T^*}\} \cup \{\mathbf{P}_\alpha : \alpha < \xi\}$ is a structure of the form

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where

- $T \subset \{\beta < \alpha : \text{Lim}(\beta)\} \times \mathcal{L}(\mathbf{aa}^+)$, and for all limit ordinals $\beta < \alpha$, $T_\beta = \{\varphi(\vec{a}) : (\beta, \varphi(\vec{a})) \in T\}$ is an $\mathcal{L}(\mathbf{aa}^+)$ -theory in vocabulary τ_0^- with parameters from J_β^T .

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- For each $\varphi \in \mathcal{L}(aa^+)$ in vocabulary τ_0^- and each $\vec{a} \in J_\beta^T$, T_β contains either $\varphi(\vec{a})$ or $\neg\varphi(\vec{a})$.
- T_β is closed under the rules for first-order logic and the axioms $(A0^+) - (A5^+)$, and weakly consistent for good formulas relative to J_β^T .
- T_β attempts to describe what $\mathcal{L}(aa^+)$ -truth in the relevant next admissibles of $(J_\beta, \in, T \restriction \beta)$, where $T \restriction \beta =_{\text{def}} T \restriction J_\beta^T$, but may be wrong about the aa^+ -quantifiers.

- If φ is a first-order formula in vocabulary τ_0^- and $\vec{a} \in J_\beta^T$, then $\varphi(\vec{a}) \in T_\beta$ if and only if $(J_\beta^T, \in, T \restriction \beta) \models \varphi(\vec{a})$.
- Conditions which describe in which way the formulas T_β are about the next admissibles. For example, the following sentence, which intuitively says that J_α^T is a member of the relative next admissible, is always in T_β :

$$\begin{aligned} & \text{aa}^+ s_1 \dots \text{aa}^+ s_m \exists y \\ & [\exists \beta \in \text{Ord} (\Psi(y, \beta, J_{\sup T}^T, s_1, \dots, s_m) \wedge \\ & \forall x (\exists \gamma \in \text{Ord} \Psi(x, \gamma, J_{\sup T}^T, s_1, \dots, s_m) \rightarrow (x \in J_{\sup T}^T \leftrightarrow x \in y)))] \end{aligned}$$

- ...

- $T^* \subset \mathcal{L}(\text{aa}^+) \times J_\alpha^T$ is an $\mathcal{L}(\text{aa}^+)$ -theory in the vocabulary τ_ξ^- with parameters from J_α^T .
- T^* is complete for good formulas in vocabulary τ_ξ^- with parameters in J_α^T , closed under first-order axioms and the axioms (A0⁺) - (A5⁺), and weakly consistent for good formulas in vocabulary τ_ξ^- .
- T^* is right about first order truth in the next admissible sets built with the predicates P_η but may be wrong about the aa⁺-quantifier.
- If a first-order formula φ is good, the **P**-predicates appearing in φ are among $\mathbf{P}_{\eta_1}, \dots, \mathbf{P}_{\eta_n}$, and $\vec{a} \in J_\alpha^T$, then

$$\varphi(\vec{a}) \in T^* \Leftrightarrow (J_\alpha^T, \in, T, P_{\eta_1}, \dots, P_{\eta_n})^+ \models \varphi(\vec{a}).$$

Some other conditions

- For any good $\mathcal{L}(\text{aa}^+)$ -formula φ and any $\vec{a} \in J_\alpha^T$, $\varphi(\vec{a})$ is in T^* if and only if $\text{aa}^+ \vec{s} \varphi(\vec{a})$ is in T^* , where none of the subset variables in $\vec{s}$ occur in φ .
- For any good $\mathcal{L}(\text{aa}^+)$ -formula φ , $\exists x (x \in J_{\sup T}^T \wedge \varphi(x, \vec{a}))$ is in T^* if and only if there is some $b \in J_\alpha^T$ such that $\varphi(b, \vec{a})$ is in T^* .
- Coherence between T and T^* for formulas not containing any **P**-predicates: For all limit $\beta < \alpha$, all $\vec{a} \in J_\beta^T$, and all very good φ in vocabulary τ_0^- ,

$$\varphi(\vec{a}) \in T_\beta \text{ if and only if } (\varphi(\vec{a})^{(J_\beta^T)})' \in T^*,$$

where $\varphi(\vec{a})^{(J_\beta^T)}$ is obtained by replacing each occurrence of $x \in J_{\sup T}^T$ by $x \in J_\beta^T$, and if $\text{aa}^+ s$ appears in φ , each $x \in s$ is replaced by $x \in s \wedge x \in J_\beta^T$.

Definition

Suppose $\mathbf{M} = (M, \in, T, T^*, (P)_\xi)$ is τ_ξ -structure and $\mathbf{N} = (N, \in, \bar{T}, \bar{T}^*, (\bar{P})_\nu)$ is a τ_ν -structure with $\xi \leq \nu$. A function $\pi : M \rightarrow N$ is called an aa⁺-embedding if it satisfies the following conditions:

1. π is a first-order elementary embedding from $(M, \in, T, (P)_\xi)$ to $(N, \in, \bar{T}, (\bar{P})_\xi)$.
2. For all good $\varphi \in \mathcal{L}(\text{aa}^+)$ in vocabulary τ_ξ^- , and $\vec{a} \in J_\alpha^T$, $\varphi(\vec{a}) \in T^*$ if and only if $\varphi(\pi(\vec{a})) \in \bar{T}^*$.

aa^+ -ultrapower

- The aa^+ -ultrapower of an aa^+ -mouse $(J_\alpha^T, \in, T, T^*)$ will be a structure of the form $(M, E, S, S^*, (P')_{\xi+1})$
- One new predicate P'_ξ which is the image $j[J_\alpha^T]$ of the domain of the mouse under the aa^+ -ultrapower embedding j

aa⁺-ultrapower

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- One new predicate P'_ξ which is the image $j[J_\alpha^T]$ of the domain of the mouse under the aa⁺-ultrapower embedding j
- For each finite $d = \{\eta_1, \dots, \eta_n\} \subset \xi + 1$, we form an auxiliary model $\mathbf{M}_d = (M_d, E_d, S_d, S_d^*, P_{\eta_1}^d, \dots, P_{\eta_n}^d)$.
- The idea is that if the ultrapower $(M, E, S, S^*, (P')_{\xi+1})$ and the auxiliary model $\mathbf{M}_d$ are well-founded, then the transitive collapse of the domain M_d is the next admissible of the collapse of $(M, E, S, S^*, P'_{\eta_1}, \dots, P'_{\eta_n})$.
- This allows us to show that if the aa⁺-ultrapower and all the auxiliary models $\mathbf{M}_d$ are well-founded, then the ultrapower collapses to an aa⁺-premouse

Definition

Suppose $\mathbf{J}_\alpha^T = (J_\alpha^T, \in, T, T^*, (P)_\xi)$ is an aa^+ -premouse.

1. $M' = M'_\emptyset$ is the set of all $\varphi(s, x, \vec{a})$ in vocabulary τ_ξ^- , where $\vec{a} \in J_\alpha^T$, such that $aa^+s \varphi(s, x, \vec{y})$ is good, and $aa^+s \exists x (x \in J_{\sup T}^T \wedge \varphi(s, x, \vec{a})) \in T^*$.

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aa⁺s $\exists x (x \in J_{\sup T}^T \wedge \varphi(s, x, \vec{a})) \in T^*$.
2. For nonempty $d \in [\xi + 1]^{<\omega}$, there are two cases. If $\xi \notin d$ and $d = \{\eta_1, \dots, \eta_n\}$, M'_d is the set of all $\varphi(s, x, \vec{a})$ in vocabulary τ_ξ^- , where $\vec{a} \in J_\alpha^T$, such that aa⁺s $\varphi(s, x, \vec{y})$ is good, and

$$\text{aa}^+s \exists x \exists \beta \in \text{Ord} (\Gamma(x, \beta, J_{\sup T}^T, \eta_1, \dots, \eta_n) \wedge \varphi(s, x, \vec{a})) \in T^*.$$

If $\xi \in d$ and $d = \{\eta_1, \dots, \eta_n, \xi\}$, then M'_d is the set of all $\varphi(s, x, \vec{a})$ in vocabulary τ_ξ^- , where $\vec{a} \in J_\alpha^T$, such that aa⁺s $\varphi(s, x, \vec{y})$ is good, and

$$\text{aa}^+s \exists x \exists \beta \in \text{Ord} (\Gamma(x, \beta, J_{\sup T}^T, \eta_1, \dots, \eta_n, s) \wedge \varphi(s, x, \vec{a})) \in T^*.$$

Definition

The equivalence relations for members of M' and M'_d are defined as follows: For $\varphi(s, x, \vec{a}), \varphi'(s, x, \vec{a}') \in M'$,

$$\varphi(s, x, \vec{a}) \sim \varphi'(s, x, \vec{a}') \text{ if } \text{aa}^+ s (f_{\varphi(s, x, \vec{a})}(s) = f_{\varphi'(s, x, \vec{a}')} (s)) \in T^*.$$

For $\varphi(s, x, \vec{a}), \varphi'(s, x, \vec{a}') \in M'_d$, suppose $e, e' \subset d$ are minimal such that $\varphi(s, x, \vec{a}) \in M'_e$ and $\varphi'(s, x, \vec{a}') \in M'_{e'}$. We define

$$\begin{aligned} \varphi(s, x, \vec{a}) \sim_d \varphi'(s, x, \vec{a}') \text{ if} \\ e = e' \text{ and } \text{aa}^+ s (f_{\varphi(s, x, \vec{a})}^{e, \xi}(s) = f_{\varphi'(s, x, \vec{a}')}^{e', \xi}(s)) \in T^*. \end{aligned}$$

Definition

Suppose $\mathbf{J}_\alpha^T = (J_\alpha^T, \in, T, T^*, (P)_\xi)$ is an aa^+ -premouse. Its aa^+ -ultrapower is the $\tau_{\xi+1}$ -structure $\mathbf{M} = (M, E, S, S^*, (P')_{\xi+1})$ defined as follows:

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1. M is the set of equivalence classes $[\varphi(s, x, \vec{a})]$ of $\sim$ on M' .
2. $[\varphi(s, x, \vec{a})]E[\varphi'(s, x, \vec{a}')] \text{ iff } aa^+s R_\in(f_{\varphi(s, x, \vec{a})}(s), f_{\varphi'(s, x, \vec{a}')} (s)) \in T^*.$

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3. $([\varphi(s, x, \vec{a})], [\varphi'(s, x, \vec{a}')]) \in S \text{ iff } \text{aa}^+ s R_T(f_{\varphi(s, x, \vec{a})}(s), f_{\varphi'(s, x, \vec{a}')} (s)) \in T^*.$
4. S^* consists of $\varphi(\mathbf{P}_\xi, [\theta_1(s, x, \vec{a}_1)], \dots, [\theta_n(s, x, \vec{a}_n)])$, where $\varphi(s, x_1, \dots, x_n) \in \mathcal{L}(\text{aa}^+)$ is in vocabulary τ_ξ^- and $\text{aa}^+ s \varphi(s, f_{\theta_1(s, x, \vec{a}_1)}(s), \dots, f_{\theta_n(s, x, \vec{a}_n)}(s)) \in T^*.$
5. $[\varphi(s, x, \vec{a})] \in P'_\gamma \text{ iff } \text{aa}^+ s P_\gamma(f_{\varphi(s, x, \vec{a})}(s)) \in T^* \text{ for } \gamma < \xi.$
6. $P'_\xi = \{j(b) : b \in J_\alpha^T\}$, where $j : J_\alpha^T \rightarrow M$ is the canonical embedding defined by $j(b) = [x = b]$.

Auxiliary models $\mathbf{M}_d$

Definition

Suppose $\mathbf{J}_\alpha^T = \langle J_\alpha^T, \in, T, T^*, (P)_\xi \rangle$ is an aa^+ -premouse. For all $d = \{\eta_1, \dots, \eta_n\} \in [\xi + 1]^{<\omega}$, the model $\mathbf{M}_d = (M_d, E_d, S_d, S_d^*, P_{\eta_1}^d, \dots, P_{\eta_n}^d)$ is defined as follows:

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- $[\varphi_1(s, x, \vec{a}_1)]_d E_d [\varphi_2(s, x, \vec{a}_2)]_d$ iff $\text{aa}^+ s R_{\in} (f_{\varphi_1(s, x, \vec{a}_1)}^{d_1}(s), f_{\varphi_2(s, x, \vec{a}_2)}^{d_2}(s)) \in T^*$, where $d_1, d_2 \subset d$ are minimal such that $\varphi_1(s, x, \vec{a}_1) \in M_{d_1}$ and $\varphi_2(s, x, \vec{a}_2) \in M_{d_2}$.
- $([\varphi_1(s, x, \vec{a}_1)]_d, [\varphi_2(s, x, \vec{a}_2)]_d) \in S_d$ iff $\text{aa}^+ s R_T (f_{\varphi_1(s, x, \vec{a}_1)}^{d_1}(s), f_{\varphi_2(s, x, \vec{a}_2)}^{d_2}(s)) \in T^*$, where $d_1, d_2 \subset d$ are minimal such that $\varphi_1(s, x, \vec{a}_1) \in M_{d_1}$ and $\varphi_2(s, x, \vec{a}_2) \in M_{d_2}$.

Auxiliary models $\mathbf{M}_d$

- If $\xi \in d$, S_d^* consists of $\varphi(\mathbf{P}_\xi, [\theta_1(s, x, \vec{a}_1)]_d, \dots, [\theta_n(s, x, \vec{a}_n)]_d)$ in vocabulary τ_d^- such that

$$\text{aa}^+ s \varphi(s, f_{\theta_1(s, x, \vec{a}_1)}^{d_1, \xi}(s), \dots, f_{\theta_n(s, x, \vec{a}_n)}^{d_n, \xi}(s)) \in T^*,$$

where each $d_i \subset d$ is minimal such that $\theta_i(s, x, \vec{a}_i) \in M_{d_i}$.

If $\xi \notin d$, S_d^* consists of $\varphi([\theta_1(s, x, \vec{a}_1)]_d, \dots, [\theta_n(s, x, \vec{a}_n)]_d)$ in vocabulary τ_d^- such that

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where each $d_i \subset d$ is minimal such that $\theta_i(s, x, \vec{a}_i) \in M_{d_i}$.

- For all $\eta_i \in d \setminus \{\xi\}$, $[\varphi(s, x, \vec{a})]_d \in P_{\eta_i}^d$ iff $\text{aa}^+ s \mathbf{P}_{\eta_i}(f_{\varphi(s, x, \vec{a})}^{d', \xi}(s)) \in T^*$, where $d' \subset d$ is minimal such that $\varphi(s, x, \vec{a}) \in M_{d'}$.
- $P_\xi^d = \{[x = b]_d : b \in J_\alpha^T\}$ if $\xi \in d$.

Key lemma

Lemma

Suppose φ is a first-order formula.

- If $d = \{\eta_1, \dots, \eta_n\} \in [\xi + 1]^{<\omega}$ is nonempty, $\xi \notin d$, and φ is in vocabulary τ_d^- , then the following are equivalent:
 - $\mathbf{M}_d^- \models \varphi([\theta_1(s, x, \vec{a}_1)]_d, \dots, [\theta_m(s, x, \vec{a}_m)]_d)$,
 - $\mathbf{aa}^+ s \varphi_d^*(f_{\theta_1(s, x, \vec{a}_1)}^{d_1}(s), \dots, f_{\theta_m(s, x, \vec{a}_m)}^{d_m}(s)) \in T^*$,where each d_i is the minimal subset of d such that $\theta_i(s, x_i, \vec{a}_i) \in d_i$.
- If $\xi \in d = \{\eta_1, \dots, \eta_n, \xi\} \in [\xi + 1]^{<\omega+1}$, and φ is in vocabulary τ_d^- , then the following are equivalent:
 - $\mathbf{M}_d^- \models \varphi(P_\xi^d, [\theta_1(s, x, \vec{a}_1)]_d, \dots, [\theta_m(s, x, \vec{a}_m)]_d)$,
 - $\mathbf{aa}^+ s \varphi_d^*(s, f_{\theta_1(s, x, \vec{a}_1)}^d(s), \dots, f_{\theta_m(s, x, \vec{a}_m)}^d(s)) \in T^*$,where each d_i is the minimal subset of d such that $\theta_i(s, x_i, \vec{a}_i) \in d_i$.

Here, e.g., φ_d^* for $\xi \in d$ is obtained by inductively replacing each subformula $\exists x \psi(x)$ of φ by

$$\exists x (\exists \beta \in \text{Ord} \Psi(x, \beta, J_{\sup T}^T, \eta_1, \dots, \eta_n, s) \wedge \psi(x))$$

Key lemma continued

Moreover, the following are equivalent

- $(M, E, S, (P')_{\xi+1}) \models \varphi([\theta_1(s, x, \vec{a}_1)], \dots, [\theta_m(s, x, \vec{a}_m)])$,
- $\mathbf{aa}^+ s \varphi_{\emptyset}^*(f_{\theta_1(s, x, \vec{a}_1)}(s), \dots, f_{\theta_m(s, x, \vec{a}_m)}(s)) \in T^*$.

M_d next admissible of the $\mathbf{aa}^+$ -ultrapower

Corollary

Suppose $\mathbf{M}$ and $\mathbf{M}_d$, where $d = \{\eta_1, \dots, \eta_n\} \neq \emptyset$, are well-founded and $\bar{\mathbf{M}} = (J_{\beta}^{\bar{T}}, \in, \bar{T}, \bar{T}^*, (\bar{P})_{\xi+1})$ and $\bar{\mathbf{M}}_d = (\bar{M}_d, \in, \bar{T}_d, (\bar{T}_d)^*, \bar{P}_{\eta_1}^d, \dots, \bar{P}_{\eta_n}^d)$ are their transitive collapses. Then $\bar{M}_d$, the domain of $\bar{\mathbf{M}}_d$, is the next admissible set of $(J_{\beta}^{\bar{T}}, \in, \bar{T}, \bar{P}_{\eta_1}, \dots, \bar{P}_{\eta_n})$.

Moreover, for any good first-order φ in vocabulary τ_d^- and any $[\theta_1(s, x, \vec{a}_1)]_d, \dots, [\theta_m(s, x, \vec{a}_m)]_d \in M_d$, we have

$$\begin{aligned} & \varphi([\theta_1(s, x, \vec{a}_1)]_d, \dots, [\theta_m(s, x, \vec{a}_m)]_d) \in (\bar{T})_d^* \\ \Leftrightarrow & \bar{\mathbf{M}}_d^- \models \varphi([\theta_1(s, x, \vec{a}_1)]_d, \dots, [\theta_m(s, x, \vec{a}_m)]_d). \end{aligned}$$

Ultrapower is an aa^+ -premouse

Lemma

If the ultrapower $\mathbf{M} = (M, \in, S, S^, (P')_{\xi+1})$ of an aa^+ -premouse $\mathbf{J}_\alpha^T = (J_\alpha^T, \in, T, T^*, (P)_\xi)$ is well-founded, then its transitive collapse $\bar{\mathbf{M}} = (J_\beta^{\bar{T}}, \in, \bar{T}, \bar{T}^*, (\bar{P})_{\xi+1})$ is also an aa^+ -premouse. Moreover, the ultrapower embedding j (composed with the collapse) is an aa^+ -embedding.*

Iterated aa^+ -ultrapowers

Definition

Suppose $\mathbf{J}_\gamma^T = (J_\gamma^T, \in, T, T^*, (P^0)_0)$ is an aa^+ -premouse with vocabulary τ_0 and δ is an ordinal. The models $\mathbf{M}_\alpha = (M_\alpha, E_\alpha, T_\alpha, T_\alpha^*, (P^\alpha)_\alpha)$, $\alpha < \delta$, and the directed system

$$\langle (M_\alpha, E_\alpha, T_\alpha, T_\alpha^*, (P^\alpha)_\alpha), j_{\alpha\beta} : \alpha < \beta < \delta \rangle,$$

called the aa^+ -iteration of $\mathbf{J}_\alpha^T$ of length δ , are defined as follows:

- $(M_0, E_0, T_0, T_0^*, (P^0)_0) = \mathbf{J}_\gamma^T$
- The vocabulary of $(M_\alpha, E_\alpha, T_\alpha, T_\alpha^*, (P^\alpha)_\alpha)$ is τ_α .
- The $j_{\alpha\beta}$ form a commuting system of aa^+ -embeddings

$$j_{\alpha\beta} : (M_\alpha, E_\alpha, T_\alpha, T_\alpha^*, (P^\alpha)_\alpha) \rightarrow (M_\beta, E_\beta, T_\beta, T_\beta^*, (P^\beta)_\beta) \upharpoonright \tau_\alpha.$$

- For successor $\alpha + 1$ we let

$$(M_{\alpha+1}, E_{\alpha+1}, T_{\alpha+1}, T_{\alpha+1}^*, (P^{\alpha+1})_{\alpha+1}) = \text{Ult}(M_\alpha, E_\alpha, T_\alpha, T_\alpha^*, (P^\alpha)_\alpha).$$

The embedding $j_{\alpha, \alpha+1}$ is the aa^+ -ultrapower embedding.

Definition continued

- For limit ν , we let $(M_\nu, E_\nu, T_\nu, T_\nu^*, (P^\nu)_\nu)$ be the direct limit of the directed system

$$\langle (M_\alpha, E_\alpha, T_\alpha, T_\alpha^*, (P^\alpha)_\alpha), j_{\alpha\beta} : \alpha < \beta < \nu \rangle.$$

The embeddings $j_{\alpha\nu}$ for $\alpha < \nu$ are the direct limit embeddings.

The models $\mathbf{M}_\alpha$ are called iterates of $\mathbf{J}_\gamma^T$.

Iterated aa^+ -ultrapowers

Definition

Suppose α is an ordinal and $d \in [\alpha + 1]^{<\omega}$.

1. $(M_\emptyset^\alpha)'$ is the set of $\varphi(s, x, \vec{a})$, where $\vec{a} \in M_\alpha$, such that $aa^+s \varphi(s, x, \vec{y})$ is good, and $aa^+s \exists x (x \in J_{\sup T}^T \wedge \varphi(s, x, \vec{a})) \in T_\alpha^*$.
2. If $\alpha \in d = \{\eta_1, \dots, \eta_n, \alpha\}$, then $(M_d^\alpha)'$ is the set of $\varphi(s, x, \vec{a})$, where $\vec{a} \in M_\alpha$, such that $aa^+ \varphi(s, x, \vec{a}) \in (M_e^\alpha)'$ for some $e \subsetneq d$ or

$$aa^+s \exists x \exists \beta \in \text{Ord} (\Gamma(\beta, x, \eta_1, \dots, \eta_n, s) \wedge \varphi(s, x, \vec{a})) \in T_\alpha^*.$$

3. If $\alpha \notin d = \{\eta_1, \dots, \eta_n\}$, then $(M_d^\alpha)'$, is the set of those good $\varphi(s, x, \vec{a})$, where $\vec{a} \in M_\alpha$, such that

$$aa^+s \exists x \exists \beta \in \text{Ord} (\Gamma(\beta, x, \eta_1, \dots, \eta_n) \wedge \varphi(s, x, \vec{a})) \in T_\alpha^*.$$

4. For members of $(M_d^\alpha)'$, $\sim_d^\alpha$ is defined as $\sim_d$ was defined in Definition 10, only replacing T^* with T_α^* . The equivalence class of $\varphi(s, x, \vec{a})$ under $\sim_d^\alpha$ is denoted by $[\varphi(s, x, \vec{a})]_d^\alpha$.

Iterated aa^+ -ultrapowers

Definition

Suppose $\mathbf{J}_\gamma^T = (J_\gamma^T, \in, T, T^*, (P^0)_0)$ is an aa^+ -premouse with vocabulary τ_0 , and the iteration of $\mathbf{J}_\gamma^T$ of length δ is as in Definition 17. For any $\alpha < \beta < \delta$ and nonempty finite $d \subset \alpha$, the models $\mathbf{M}_d^\alpha = (M_d^\alpha, E_d^\alpha, T_d^\alpha, (T_d^\alpha)^*, P_{d, \eta_1}^\alpha, \dots, P_{d, \eta_n}^\alpha)$, where $d = \{\eta_1, \dots, \eta_n\}$, and the maps $j_d^{\alpha\beta} : \mathbf{M}_d^\alpha \rightarrow \mathbf{M}_d^\beta$ are defined inductively as follows:

1. If $\alpha = \lambda + 1$, then $\mathbf{M}_d^\alpha$ is defined using T_λ^* and the equivalence classes $[\varphi(s, x, \vec{a})]_d^\lambda$ as $\mathbf{M}_d$ was defined in 13 using T^* and $[\varphi(s, x, \vec{a})]_d$.
2. If $\alpha = \lambda + 1$, $j_d^{\alpha, \alpha+1} (= j_d^{\lambda+1, \lambda+2})$ is defined by

$$j_d^{\alpha, \alpha+1}([\varphi(s, x, \vec{a})]_d^\lambda) = [\varphi(\mathbf{P}_\lambda, x, j_{\lambda, \alpha}(\vec{a}))]_d^\alpha.$$

3. Suppose α is a limit, and $\mathbf{M}_d^\lambda$ and $j_d^{\lambda\eta}$ have been defined for all $\lambda < \eta < \alpha$ such that $d \subset \lambda$, and $\langle \mathbf{M}_d^\lambda, j_d^{\lambda\eta} : d \subset \lambda < \eta < \alpha \rangle$ is a directed system. Then $\mathbf{M}_d^\alpha$ is a direct limit of the system $\langle \mathbf{M}_d^\gamma, j_d^{\lambda\eta} : d \subset \lambda < \eta < \alpha \rangle$ and $j_d^{\lambda\alpha}$, $\lambda < \alpha$, are the direct limit maps.

Definition continued

- Suppose α is a limit, $\mathbf{M}_d^\lambda$ has been defined for each $\lambda \leq \alpha$ such that $d \subset \lambda$, and $j_d^{\lambda\eta}$ has been defined for all $d \subset \lambda < \eta \leq \alpha$. Then $j_d^{\alpha, \alpha+1}$ is defined as follows. Suppose $y \in M_d^\alpha$ is such that $y = j_d^{\lambda+1, \alpha}([\varphi(s, x, \vec{a})]_d^\lambda)$ for some $[\varphi(s, x, \vec{a})]_d^\lambda \in M_d^{\lambda+1}$. Then we define

$$j_d^{\alpha, \alpha+1}(y) = [\varphi(\mathbf{P}_\lambda, x, j_{\lambda\alpha}(\vec{a}))]_d^\alpha.$$

Iterated aa^+ -ultrapowers

Lemma

Suppose $\langle (M_\alpha, E_\alpha, T_\alpha, T_\alpha^*, (P^\alpha)_\alpha), j_{\alpha\beta} : \alpha < \beta < \delta \rangle$ is the aa^+ -iteration of $(M_0, \in, T_0, T_0^*)$ of length δ and for all $\alpha < \delta$ and nonempty $d \in [\alpha]^{<\omega}$ the models $\mathbf{M}_d^\alpha = (M_d^\alpha, E_d^\alpha, T_d^\alpha, (T_d^\alpha)^*, P_{d,\eta_1}^\alpha, \dots, P_{d,\eta_n}^\alpha)$ are as in the preceding definition. Suppose that for all $\alpha < \delta$ and all nonempty $d \in [\alpha]^{<\omega}$, M_α and M_d^α are well-founded. Then the following hold:

- For all $\alpha < \delta$, $\mathbf{M}_\alpha$ is (isomorphic to) an aa^+ -premouse.
- For all $\alpha < \beta < \delta$, $j_{\alpha\beta}$ is an aa^+ -embedding.
- For all $\alpha < \delta$ and $d = \{\eta_1, \dots, \eta_n\} \in [\alpha]^{<\omega}$, M_d^α is the next admissible of $(M_\alpha, \in, T_\alpha, P_{\eta_1}^\alpha, \dots, P_{\eta_n}^\alpha)$.
- For all $\alpha < \delta$, $d \in [\alpha]^{<\omega}$, all good first-order φ in vocabulary τ_d^- , and all $\vec{a} \in M_d^\alpha$, $\varphi(\vec{a}) \in (T_d^\alpha)^*$ if and only if $(\mathbf{M}_d^\alpha)^- \models \varphi(\vec{a})$.
- For all $\alpha < \beta < \delta$, $j_d^{\alpha\beta}$ is an aa^+ -embedding.

Lemma

Suppose $\langle (M_\alpha, E_\alpha, T_\alpha, T_\alpha^*, (P^\alpha)_\alpha), j_{\alpha\beta} : \alpha < \beta < \delta \rangle$ is the aa^+ -iteration of an aa^+ -mouse $(M_0, \in, T_0, T_0^*)$ of length δ , and for $d \subset \alpha < \beta < \delta$ the models $\mathbf{M}_d^\alpha = (M_d^\alpha, E_d^\alpha, T_d^\alpha, (T_d^\alpha)^*, P_{d, \eta_1}^\alpha, \dots, P_{d, \eta_n}^\alpha)$ and maps $j_d^{\alpha\beta} : \mathbf{M}_d^\alpha \rightarrow \mathbf{M}_d^\beta$, are as in Definition 19. Suppose further that κ is a regular cardinal such that $|M_0| < \kappa$ and $\kappa < \delta$.

Then for any $d \in [\kappa]^{<\omega}$, any good formula φ in vocabulary τ_d^- , and any $\vec{a} \in M_d^\kappa$,

$$\varphi(\vec{a}) \in (T_d^\kappa)^* \Leftrightarrow ((\mathbf{M}_d^\kappa)^-; M_\kappa) \models^+ \varphi_\kappa(\vec{a}),$$

where φ_κ is obtained from φ by replacing each quantifier aa^+s by aa_κ^+s .

The main result

Theorem

Suppose $\langle (M_\alpha, E_\alpha, T_\alpha, T_\alpha^, (P^\alpha)_\alpha), j_{\alpha\beta} : \alpha < \beta < \delta \rangle$ is the aa^+ -iteration of an aa^+ -mouse $(M_0, \in, T_0, T_0^*)$ of length δ . If $\kappa < \delta$ is a regular cardinal such that $|M_0| < \kappa$, then the domain $M_\kappa = J_{\beta_\kappa}^{T_\kappa}$ of the κ -iterate is the level $J_{\beta_\kappa}^{\text{Tr}^\kappa}$ of $C(\text{aa}^+)$. In particular, if M_0 is countable, then M_{ω_1} is the level $J_{\beta_{\omega_1}}^{\text{Tr}}$ of $C(\text{aa}^+)$.*

Thank you!

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